

*The second test for this course will be given in class on Wednesday, April 14. The test covers all of the material that we have studied from Chapter 3 and Sections 4.1 through 4.3 in Chapter 4, except for Sections 3.3 and 3.7. (This means that you will not be asked about `grep`, or about other Linux commands, or about the complex syntax that is used for practical regular expressions. And there will be no Pumping Lemma proof. Note that pushdown automata, Section 4.4, are not included.)*

*The test will include some “short essay” questions that ask you to define something, or discuss something, or explain something, and so on. Other than that, you can expect most of the questions to be similar to problems that have been given on the homework. There might be some simple proofs.*

### **Here are some terms and ideas that you should be familiar with for the test:**

alphabet (finite, non-empty set of “symbols”)  
 string over an alphabet  $\Sigma$   
 length of a string,  $|x|$   
 empty string,  $\varepsilon$   
 concatenation of strings,  $xy$  or  $x \cdot y$   
 reverse of a string,  $x^R$   
 $x^n$ , for a string  $x$  and a natural number  $n$   
 $n_\sigma(x)$ , the number of occurrences of a symbol  $\sigma$  in a string  $x$   
 the set of all possible strings over  $\Sigma$ , denoted  $\Sigma^*$   
 language over an alphabet  $\Sigma$  (a subset of  $\Sigma^*$ )  
 a language over  $\Sigma$  is an element of  $\mathcal{P}(\Sigma^*)$   
 the set of strings over  $\Sigma$  is countable; the set of languages over  $\Sigma$  is uncountable  
 union, intersection, set difference, and complement applied to languages  
 concatenation of languages:  $LM, L^n$  for  $n \in \mathbb{N}$   
 Kleene-star operation on a language:  $L^*$   
 regular expression over an alphabet  $\Sigma$ ; the operators:  $*$ ,  $|$ , and concatenation  
 regular language; the language  $L(r)$  generated by a regular expression  $r$   
 DFA (Deterministic Finite Automaton)  
 transition diagram [the usual picture] of a DFA  
 state (in a finite-state automaton); start state; accepting state (also known as final state)  
 definition of a DFA as a list of five things,  $(Q, \Sigma, q_o, \delta, F)$  — and what each thing means  
 how a DFA computes (that is, what it does when it reads and processes a string)  
 NFA (Non-deterministic Finite Automaton); the differences between NFAs and DFAs  
 nondeterminism  
 $\varepsilon$ -transitions  
 what it means for an NFA to accept a string  
 the language,  $L(M)$ , accepted by an NFA or DFA,  $M$   
 algorithm for converting an NFA to an equivalent DFA

algorithm for converting a regular expression to an equivalent NFA  
DFAs, NFAs, and regular expressions all define the same class of languages  
operations ( $L \cup M$ ,  $L \cap M$ ,  $LM$ ,  $L^*$ ,  $\bar{L}$ ,  $L^R$ ) on regular languages produce regular languages  
CFGs (Context-Free Grammars)

production rules; non-terminal and terminal symbols; start symbol

definition of a CFG as a list of 4 things,  $G = (V, \Sigma, P, S)$

derivation (of a string from the start symbol of a CFG)

the language,  $L(G)$ , generated by a CFG,  $G$ ; context-free language

if  $L$  and  $M$  are context-free languages, then so are  $L \cup M$ ,  $LM$ , and  $L^*$

every regular language is context-free

BNF (Backus-Naur Form)

using BNF to define the syntax of a language

parsing

parse tree

left derivations and right derivations

how a parse tree corresponds to a derivation

examples of languages that are not regular, such as:

$$\{a^n b^n \mid n = m\}$$

$$\{a^n b^m c^k \mid k = n + m\}$$

$$\{w \in \{a, b\}^* \mid w = w^R\}$$

$$\{w \in \{a, b\}^* \mid n_a(w) < n_b(w)\}$$

$$\{ww \mid w \in \{a, b\}^*\}$$

$$\{a^n b^n c^n \mid n \in \mathbb{N}\}$$

$$\{a^{n^2} \mid n \in \mathbb{N}\}$$

$$\{www \mid w \in \{a, b\}^*\}$$

examples of languages that are context-free but not regular, such as:

the **first four** languages in the previous list

some of the tasks that you could be asked to perform:

finding a regular expression for a given language

finding a DFA or NFA for a given language

finding a regular expression for an NFA or DFA

converting an NFA into an equivalent DFA, using the algorithm

converting a regular expression into an equivalent NFA, using the algorithm

determining whether a given string is accepted by an NFA or DFA

finding a derivation for a given string from a given CFG

finding a CFG for a given language

finding the language generated by a given CFG

using BNF in basic ways

finding a parse tree for a given string from a given CFG

finding a left derivation or a right derivation

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Here are some practice problems from old exams...

Note that this is **not** meant to be a complete review of everything that might be on the test. It is meant to give you a feeling for what kinds of test questions are possible

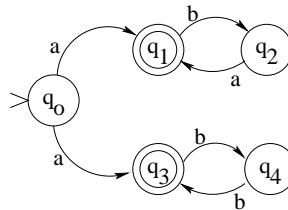
1. a) Suppose that  $L$  and  $M$  are languages over an alphabet  $\Sigma$ . How is the language  $LM$  defined?  
 b) Suppose that  $\Sigma = \{a, b, c\}$  and that  $L$  is the language  $L = \{a, b\}$  and  $M$  is the language  $M = \{w \in \Sigma^* \mid w \text{ ends with a } c\}$ . What strings are in the language  $LM$ ? (Describe the language with words, **not** with a regular expression.)

2. Consider the following context-free grammars:

$S \rightarrow TR$	$S \rightarrow SS$
$T \rightarrow aTb$	$S \rightarrow T$
$T \rightarrow aT$	$T \rightarrow aTb$
$T \rightarrow a$	$T \rightarrow \varepsilon$
$R \rightarrow aRb$	
$R \rightarrow bR$	
$R \rightarrow b$	

- a) For the grammar on the **left**, find a *left derivation* and the corresponding *parse tree* for the string  $aaabbabbb$ .
- b) For the grammar on the **right**, find **two** *left derivations* and the two corresponding *parse trees* for the string  $abaabbab$ .

3. Consider the following Nondeterministic Finite Automaton:



- a) Give a regular expression for the language accepted by this NFA.
  - b) Apply the NFA-to-DFA conversion algorithm to construct a DFA that accepts the same language as this NFA.
4. Consider the regular language  $L = \{w \in \{a, b\}^* \mid w \text{ contains the substring } abaab\}$ . Draw a DFA that accepts exactly this language.

5. Give a regular expression that generates each language (no explanation necessary):

$$L_1 = \{w \in \{a, b\}^* \mid w \text{ contains at least 2 } a\text{'s}\}$$

$$L_2 = \{w \in \{a, b\}^* \mid w \text{ contains exactly 2 } a\text{'s}\}$$

6. Draw an NFA that accepts the language over the alphabet  $\Sigma = \{a, b, c\}$  that is generated by the regular expression  $(a|b)^*cc^*(a|b)^*$ . You do **not** have to use a specific regular-expression-to-NFA conversion algorithm; any NFA that works will do.
7. Let  $L$  be a language over some alphabet  $\Sigma$ . Suppose that  $\varepsilon \in L$ . Show that  $L \subseteq L^2$ .

8. What does it mean that  $x^R = x$  for a string  $x$ ? Let  $\Sigma = \{a, b, c\}$ . Give several examples of strings,  $x$ , over the alphabet  $\Sigma$  that have the property that  $x^R = x$ , and explain in English what this property means.
9. Give a Context-Free Grammar for the language  $L = \{a^n b^m c b^m a^n \mid n \in \mathbb{N}, m \in \mathbb{N}\}$ , and briefly explain how your grammar works.
10. (“ $a^n b^m$  variations”) Match each language with the grammar that generates it. Enter the number for the correct grammar. No explanation is necessary.

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|--|---|--|
| <p>a) <math>\{a^n b^m \mid n \neq m\}</math><br/>Grammar # _____</p> | <p>b) <math>\{a^n b^m \mid n \geq m\}</math><br/>Grammar # _____</p>        | <p>c) <math>\{a^n b^m \mid n &gt; 2m\}</math><br/>Grammar # _____</p>  |
| <p>d) <math>\{a^n b^m \mid n &gt; 0\}</math><br/>Grammar # _____</p> | <p>e) <math>\{a^n b^m \mid n &gt; m &gt; 0\}</math><br/>Grammar # _____</p> | <p>f) <math>\{a^n b^m \mid n = 2m + 1\}</math><br/>Grammar # _____</p> |

1)  $S \rightarrow aSB$   
 $S \rightarrow \varepsilon$   
 $B \rightarrow b$   
 $B \rightarrow \varepsilon$

2)  $S \rightarrow aS$   
 $S \rightarrow Sb$   
 $S \rightarrow a$

3)  $S \rightarrow aaSb$   
 $S \rightarrow aT$   
 $T \rightarrow aT$   
 $T \rightarrow \varepsilon$

4)  $S \rightarrow aaTb$   
 $T \rightarrow aTb$   
 $T \rightarrow aT$   
 $T \rightarrow \varepsilon$

5)  $S \rightarrow aSb$   
 $S \rightarrow T$   
 $S \rightarrow R$   
 $T \rightarrow aT$   
 $R \rightarrow Rb$   
 $T \rightarrow a$   
 $R \rightarrow b$

6)  $S \rightarrow aaSb$   
 $S \rightarrow a$

11. Let  $G = (V, \Sigma, P, S)$  be a Context-Free Grammar.  $L(G)$  is the language generated by  $G$ . How is  $L(G)$  defined? That is, exactly what does it mean for a string,  $w$ , to be in  $L(G)$ ? (Mention  $\Sigma$ ,  $P$ , and  $S$  in your answer.)
12. NFA stands for “Nondeterministic Finite Automaton.” Discuss how an NFA works and what it means to say that it is “nondeterministic.” (What does an NFA do as it reads an input string? What does it mean to say that the NFA accepts the string?)